

# Contents

<b>1</b>	<b>zerocash notes, and highlights</b>	<b>1</b>
1.1	zerocoin drawbacks . . . . .	1
<b>2</b>	<b>minting</b>	<b>1</b>
<b>3</b>	<b>spending</b>	<b>1</b>
3.0.1	pouring . . . . .	1

## 1 zerocash notes, and highlights

### 1.1 zerocoin drawbacks

zerocoin isn't used in daily transaction due to performance limitations, but rather for anonymization, or laundering coins through decentralized mix. performance bottleneck is that redeeming zerocoins requires double-discrete-logarithm proofs of knowledge which have size that exceeds 45kb, and 450ms to verify(at the 128-bit security level), it uses fixed denominations, can't pay in zerocoin directly, provides anonymity for the original address only.

## 2 minting

minting coin  $c := ((a_{pk}, PK_{enc}), v, \rho, r, s, cm)$  is:

$$tx_{mint} := (v, k, s, cm)$$

$$cm := COMM_s(v||k)$$

$$a_{pk} := PRF_{a_{sk}}^{addr}(0)$$

$$k := COMM_r(a_{pk}||\rho)$$

## 3 spending

Spending the coin  $c$ :

$$tx_{spend} := (cm, sn, \pi)$$

$$sn := PRF_{a_{sk}}^{sn}(\rho)$$

### 3.0.1 pouring

pouring  $coin^{old}$  into  $coin_1^{new}, coin_2^{new}$ . with  $v^{old} = v_{pub} + v_1^{new} + v_2^{new}$  as follows:

$$tx_{pour} := (rt, sn^{old}, cm_1^{new}, cm_2^{new}, \pi_{pour}, enc_{pk_{enc,1}^{new}}(C_1), enc_{pk_{enc,2}^{new}}(C_2))$$

$$C_i = (v_i^{new}, \rho_i^{new}, r_i^{new}, s_i^{new})$$

$\pi_{pour}$  is the pouring sk-snark proof of the spending/pouring process. address for each participant is the pair  $(addr_{pk}, addr_{sk})$ ,  $addr_{pk} = (a_{enc}, pk_{enc})$ ,  $addr_{sk} = (a_{sk}, sk_{enc})$